

# Finite Frames and Optimal Subspace Packings

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Party like it's 1999...



# A Research Problem

In the fall of 2000, inspired by a talk by Ed Saff at a conference in Bommerholz and a follow-up question by Hans Feichtinger, John asked me the following question (paraphrased):

*How is the problem of equally-distributing points on a sphere related to finite unit norm tight frames?*

This talk is the 2019 progress update.

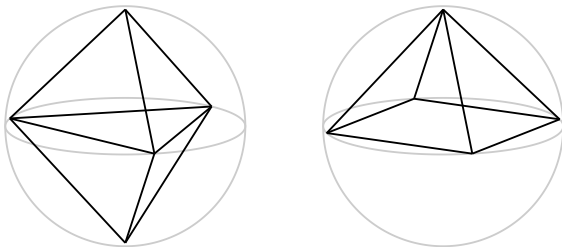
# **Optimal Packings on Spheres**

# Spherical equidistribution: Thomson vs. Tammes

Over all sets of  $N$  unit vectors  $\{\mathbf{x}_n\}_{n=1}^N$  in  $\mathbb{R}^D$ , we can try to:

- ▶ minimize  $\sum_{n=1}^N \sum_{\substack{n'=1 \\ n' \neq n}}^N \frac{1}{\|\mathbf{x}_n - \mathbf{x}_{n'}\|}$  (**Thomson, 1904**)
- ▶ maximize  $\min_{n \neq n'} \|\mathbf{x}_n - \mathbf{x}_{n'}\|$  (**Tammes, 1930**)

For example, when  $N = 5$ ,  $D = 3$ :



# Solving Tammes in the Simplest Case

**Theorem:** [Rankin 55] When  $N \leq D + 1$ , every solution to Tammes problem is a  $N$ -vector regular simplex.

**Proof:** For any unit vectors  $\{\mathbf{x}_n\}_{n=1}^N$  in  $\mathbb{R}^D$ ,

$$\|\mathbf{x}_n - \mathbf{x}_{n'}\|^2 = 2(1 - \langle \mathbf{x}_n, \mathbf{x}_{n'} \rangle).$$

Thus,  $\operatorname{argmax}_{\{\mathbf{x}_n\}} \min_{n \neq n'} \|\mathbf{x}_n - \mathbf{x}_{n'}\| = \operatorname{argmin}_{\{\mathbf{x}_n\}} \max_{n \neq n'} \langle \mathbf{x}_n, \mathbf{x}_{n'} \rangle$ . Also,

$$0 \leq \left\| \sum_{n=1}^N \mathbf{x}_n \right\|^2 = \sum_{n=1}^N \sum_{n'=1}^N \langle \mathbf{x}_n, \mathbf{x}_{n'} \rangle \leq N + N(N-1) \max_{n \neq n'} \langle \mathbf{x}_n, \mathbf{x}_{n'} \rangle.$$

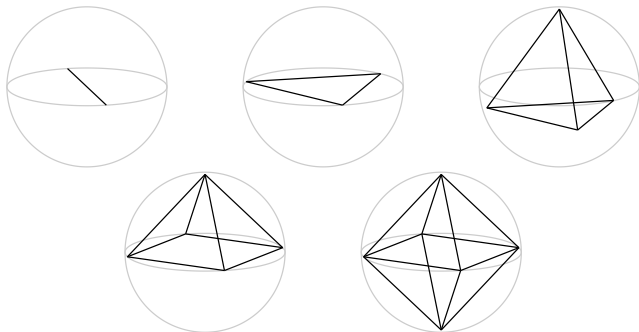
Equality only holds  $\Leftrightarrow \sum_{n=1}^N \mathbf{x}_n = \mathbf{0}$  and  $\langle \mathbf{x}_n, \mathbf{x}_{n'} \rangle$  is constant over all  $n \neq n'$ .

# Solving Tammes in the Next Simplest Case

**Theorem:** [Rankin 55]  $\max_{n \neq n'} \langle \mathbf{x}_n, \mathbf{x}_{n'} \rangle \geq 0$  when  $N \geq D + 2$ .

Moreover, for  $N \leq 2D$ , this bound can be achieved.

**Example:**  $D = 3$ ,  $N = 2, 3, 4, 5, 6$ :



# Finite Unit-Norm Tight Frames



# Notation

Let  $\mathbb{F}$  be either  $\mathbb{R}$  or  $\mathbb{C}$ . We usually regard  $N$  vectors  $\{\varphi_n\}_{n=1}^N$  in  $\mathbb{F}^D$  as the columns of a  $D \times N$  matrix

$$\Phi = [\varphi_1 \ \dots \ \varphi_N].$$

Multiplying  $\Phi$  by its  $N \times D$  conjugate-transpose  $\Phi^*$  gives its

$$\blacktriangleright N \times N \text{ Gram matrix } \Phi^* \Phi = \begin{bmatrix} \langle \varphi_1, \varphi_1 \rangle & \cdots & \langle \varphi_1, \varphi_N \rangle \\ \vdots & \ddots & \vdots \\ \langle \varphi_N, \varphi_1 \rangle & \cdots & \langle \varphi_N, \varphi_N \rangle \end{bmatrix}$$

$$\blacktriangleright D \times D \text{ frame operator } \Phi \Phi^* = \sum_{n=1}^N \varphi_n \varphi_n^*$$

In this talk, every  $\varphi_n$  is unit-norm, meaning the diagonal of  $\Phi^* \Phi$  is all ones and  $\Phi \Phi^*$  is a sum of rank-one projections.

# Orthonormal Bases (ONBs)

**Fact:** If  $\{\varphi_n\}_{n=1}^N$  is an ONB for  $\mathbb{F}^N$  then  $\Phi$  is square and satisfies  $\Phi^* \Phi = \mathbf{I}$ . Thus,  $\Phi^* = \Phi^{-1}$  and so we also have

$$\Phi \Phi^* = \mathbf{I}, \quad \text{i.e.,} \quad \mathbf{x} = \Phi \Phi^* \mathbf{x} = \sum_{n=1}^N \langle \varphi_n, \mathbf{x} \rangle \varphi_n, \quad \forall \mathbf{x} \in \mathbb{F}^N.$$

**Example:**  $\Phi = \frac{1}{\sqrt{7}} \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & \omega & \omega^2 & \omega^3 & \omega^4 & \omega^5 & \omega^6 \\ 1 & \omega^2 & \omega^4 & \omega^6 & \omega & \omega^3 & \omega^5 \\ 1 & \omega^3 & \omega^6 & \omega^2 & \omega^5 & \omega^1 & \omega^4 \\ 1 & \omega^4 & \omega & \omega^5 & \omega^2 & \omega^6 & \omega^3 \\ 1 & \omega^5 & \omega^3 & \omega & \omega^6 & \omega^4 & \omega^2 \\ 1 & \omega^6 & \omega^5 & \omega^4 & \omega^3 & \omega^2 & \omega \end{bmatrix}, \quad \omega = \exp\left(\frac{2\pi i}{7}\right).$

# Finite Unit-Norm Tight Frames (FUNTFs)

**Definition:** Unit vectors  $\{\varphi_n\}_{n=1}^N$  in  $\mathbb{F}^D$  form a **FUNTF** for  $\mathbb{F}^D$  if there exists  $C > 0$  such that

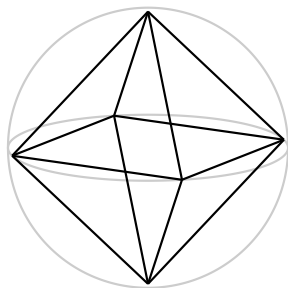
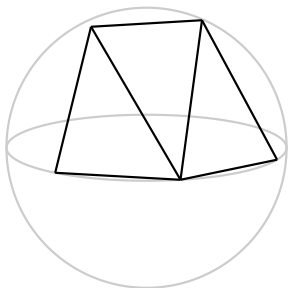
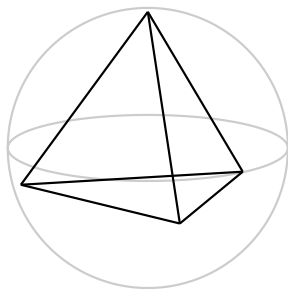
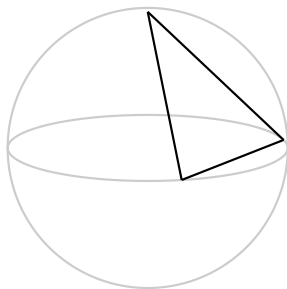
$$\Phi\Phi^* = C\mathbf{I}, \quad \text{i.e.,} \quad C\mathbf{x} = \Phi\Phi^*\mathbf{x} = \sum_{n=1}^N \langle \varphi_n, \mathbf{x} \rangle \varphi_n, \quad \forall \mathbf{x} \in \mathbb{F}^N.$$

Here,  $C = \frac{N}{D}$  since  $CD = \text{Tr}(\Phi\Phi^*) = \text{Tr}(\Phi^*\Phi) = N$ .

**Example:** Scaling the any three rows of the previous matrix gives a complex FUNTF(3, 7). For example, for rows  $\{1, 2, 4\}$ ,

$$\Phi = \frac{1}{\sqrt{3}} \begin{bmatrix} 1 & \omega & \omega^2 & \omega^3 & \omega^4 & \omega^5 & \omega^6 \\ 1 & \omega^2 & \omega^4 & \omega^6 & \omega & \omega^3 & \omega^5 \\ 1 & \omega^4 & \omega & \omega^5 & \omega^2 & \omega^6 & \omega^3 \end{bmatrix}, \quad \omega = \exp\left(\frac{2\pi i}{7}\right).$$

# Some real FUNTFs for $\mathbb{R}^3$ with $N = 3, 4, 5, 6$



**Relating FUNTFs  
to the  
Tammes Problem**

# A Big Idea from Conway, Hardin, Sloane 96

A unit vector  $\varphi$  **lifts** to a rank-one projection  $\varphi\varphi^*$ . The set

$$\{\varphi\varphi^* : \varphi \in \mathbb{F}^D, \|\varphi\| = 1\}$$

is a **projective space** and lies in the real space of all  $D \times D$  self-adjoint operators, which is a Hilbert space under the Frobenius inner product  $\langle \mathbf{A}, \mathbf{B} \rangle_{\text{Fro}} := \text{Tr}(\mathbf{A}^* \mathbf{B})$ .

Moreover, for unit vectors  $\{\varphi_n\}_{n=1}^N$  and any  $n, n'$ ,

$$\langle \varphi_n \varphi_n^*, \varphi_{n'} \varphi_{n'}^* \rangle_{\text{Fro}} = \text{Tr}(\varphi_n \varphi_n^* \varphi_{n'} \varphi_{n'}^*) = |\langle \varphi_n, \varphi_{n'} \rangle|^2,$$

and so the squared-distance between two such projections is:

$$\|\varphi_n \varphi_n^* - \varphi_{n'} \varphi_{n'}^*\|_{\text{Fro}}^2 = \text{Tr}[(\varphi_n \varphi_n^* - \varphi_{n'} \varphi_{n'}^*)^2] = 2(1 - |\langle \varphi_n, \varphi_{n'} \rangle|^2).$$

# Applying a Trivial Bound in Projective Space

**Theorem:** [Rankin 56] For any unit vectors  $\{\varphi_n\}_{n=1}^N$  in  $\mathbb{F}^D$ ,

$$\frac{N^2}{D} \leq \sum_{n=1}^N \sum_{n'=1}^N |\langle \varphi_n, \varphi_{n'} \rangle|^2$$

where equality holds if and only if  $\{\varphi_n\}_{n=1}^N$  is a FUNTF for  $\mathbb{F}^D$ .

**Proof:**

$$\begin{aligned} 0 &\leq \left\| \sum_{n=1}^N (\varphi_n \varphi_n^* - \frac{1}{D} \mathbf{I}) \right\|_{\text{Fro}}^2 = \text{Tr}[(\Phi \Phi^* - \frac{N}{D} \mathbf{I})^2] \\ &= \sum_{n=1}^N \sum_{n'=1}^N |\langle \varphi_n, \varphi_{n'} \rangle|^2 - \frac{N^2}{D}. \end{aligned}$$

# FUNTF Characterization and Construction

**Theorem:** [Benedetto, F 03] When  $N \geq D$ , every *local* minimizer of the **frame potential**

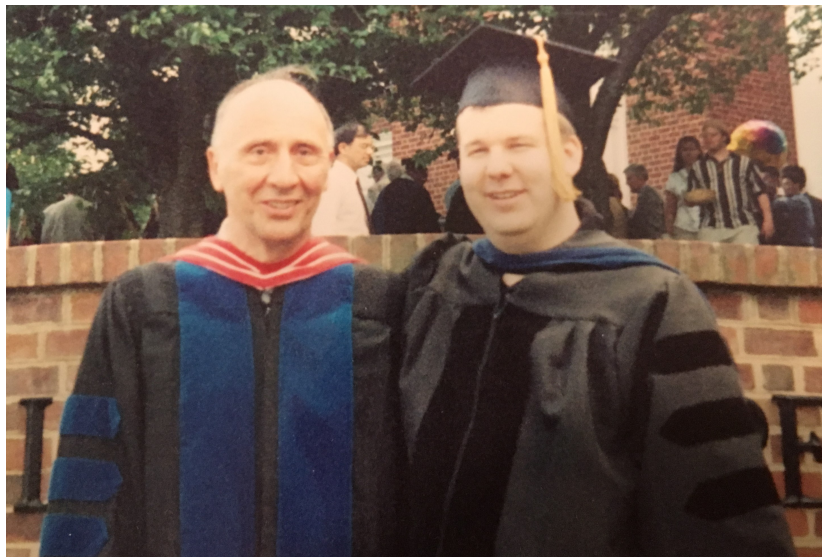
$$\sum_{n=1}^N \sum_{n'=1}^N |\langle \varphi_n, \varphi_{n'} \rangle|^2$$

is a FUNTF (and so is necessarily a global minimizer).

**Theorem:** [Cahill, F, Mixon, Poteet, Strawn 13] *Every* FUNTF can be explicitly constructed from **eigensteps**.



# Born Again



# Equiangular Tight Frames (ETFs)

**Theorem:** [Strohmer, Heath 03]

Any unit vectors  $\{\varphi_n\}_{n=1}^N$  in  $\mathbb{F}^D$  satisfy the **Welch bound**:

$$\max_{n \neq n'} |\langle \varphi_n, \varphi_{n'} \rangle| \geq \sqrt{\frac{N-D}{D(N-1)}},$$

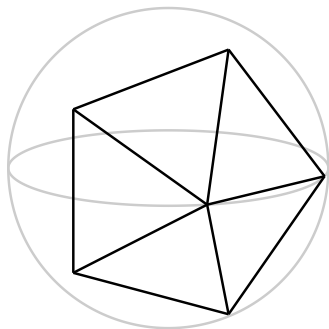
and achieve equality  $\Leftrightarrow \{\varphi_n\}_{n=1}^N$  is an ETF for  $\mathbb{F}^D$ , namely a FUNTF where  $|\langle \varphi_n, \varphi_{n'} \rangle|$  is constant over all  $n \neq n'$ .

**Proof:** Apply Rankin's simplex bound to  $\{\varphi_n \varphi_n^* - \frac{1}{D} \mathbf{I}\}_{n=1}^N$ :

$$\frac{N^2}{D} \leq \sum_{n=1}^N \sum_{n'=1}^N |\langle \varphi_n, \varphi_{n'} \rangle|^2 \leq N + N(N-1) \max_{n \neq n'} |\langle \varphi_n, \varphi_{n'} \rangle|^2.$$

See also: Rankin 56; Welch 74; Conway, Hardin, Sloane 96].

## Example: A 6-vector ETF for $\mathbb{R}^3$



$$\Phi\Phi^* = \begin{bmatrix} 2 & 0 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 2 \end{bmatrix}, \quad \Phi^*\Phi = \mathbf{I} + \frac{1}{\sqrt{5}} \begin{bmatrix} 0 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & -1 & -1 & 1 \\ 1 & 1 & 0 & 1 & -1 & -1 \\ 1 & -1 & 1 & 0 & 1 & -1 \\ 1 & -1 & -1 & 1 & 0 & 1 \\ 1 & 1 & -1 & -1 & 1 & 0 \end{bmatrix}$$

# Some Remarks on the (Rankin-)Welch Bound

- ▶ Following [Rankin 55], Rankin studied packing *antipodal pairs* of points of spheres and discovered the Welch bound about two decades before Welch [Rankin 56].
- ▶ The Welch bound is equivalent to

$$\max_{n \neq n'} \|\varphi_n \varphi_n^* - \varphi_{n'} \varphi_{n'}^*\|_{\text{Fro}}^2 \leq \frac{2N(D-1)}{D(N-1)}. \quad (1)$$

In particular, if an ETF( $D, N$ ) exists, then every optimal packing of  $N$  lines in  $\mathbb{F}^D$  is necessarily tight.

- ▶ [Conway, Hardin, Sloane 96] calls (1) the **simplex bound** since it's achieved  $\Leftrightarrow \{\varphi_n \varphi_n^* - \frac{1}{D} \mathbf{I}\}_{n=1}^N$  is a simplex. They also consider subspaces of dimension  $> 1$ .

## More Remarks on the (Rankin-)Welch Bound

- ▶ **(Gerzon)** If  $\{\varphi_n \varphi_n^* - \frac{1}{D} \mathbf{I}\}_{n=1}^N$  is a simplex, then

$$N \leq \frac{D(D+1)}{2} \text{ when } \mathbb{F} = \mathbb{R}, \quad N \leq D^2 \text{ when } \mathbb{F} = \mathbb{C}.$$

- ▶ For larger  $N$ , applying Rankin's other bound to  $\{\varphi_n \varphi_n^* - \frac{1}{D} \mathbf{I}\}_{n=1}^N$  gives the **orthoplex bound**:

$$\max_{n \neq n'} |\langle \varphi_n, \varphi_{n'} \rangle| \geq \frac{1}{\sqrt{D}}.$$

- ▶ An ETF with  $N = D^2$  is a **SIC-POVM**. Zauner has conjectured that these exist for all  $D$  [Zauner 99].
- ▶ ETFs arise in algebraic coding theory [Grey 62], quantum information theory [Zauner 99], wireless communication [Strohmer, Heath 03], and compressed sensing.

# Equiangular Tight Frames

# Harmonic ETFs: Difference Sets

**Definition:** Extracting rows from the character table of a finite abelian group  $\mathcal{G}$  yields a **harmonic frame**.

**Example:**  $\mathcal{G} = \mathbb{Z}_7$ ,  $\mathcal{D} = \{1, 2, 4\}$ ,

$$\Phi = \frac{1}{\sqrt{3}} \begin{bmatrix} 1 & \omega & \omega^2 & \omega^3 & \omega^4 & \omega^5 & \omega^6 \\ 1 & \omega^2 & \omega^4 & \omega^6 & \omega & \omega^3 & \omega^5 \\ 1 & \omega^4 & \omega & \omega^5 & \omega^2 & \omega^6 & \omega^3 \end{bmatrix}, \quad \omega = \exp\left(\frac{2\pi i}{3}\right).$$

**Theorem:** [Turyn 65] The harmonic ETF arising from  $\mathcal{D} \subseteq \mathcal{G}$  is an ETF for  $\mathbb{C}^{\mathcal{D}} \Leftrightarrow \mathcal{D}$  is a **difference set** for  $\mathcal{G}$ .

**Idea:**

$$\varphi_n \varphi_n^* = \frac{1}{3} \begin{bmatrix} \omega^n \\ \omega^{2n} \\ \omega^{4n} \end{bmatrix} [\omega^{-n} \ \omega^{-2n} \ \omega^{-4n}] = \frac{1}{3} \begin{bmatrix} 1 & \omega^{6n} & \omega^{4n} \\ \omega^n & 1 & \omega^{5n} \\ \omega^{3n} & \omega^{2n} & 1 \end{bmatrix}.$$

# Steiner ETFs

**Theorem:** [Goethals, Seidel 70] Every **balanced incomplete block design (BIBD)** with  $\Lambda = 1$  yields an ETF.

**Example:** Combine  $\begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$  and  $\begin{bmatrix} + & - & + & - \\ + & + & - & - \\ + & - & - & + \end{bmatrix}$  to form

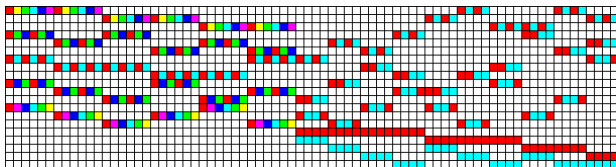
$$\Phi = \frac{1}{\sqrt{3}} \begin{bmatrix} + & - & + & - & + & - & + & - & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & + & - & + & - & + & - & + & - & + & - \\ + & + & - & - & 0 & 0 & 0 & 0 & + & + & - & - & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & + & + & - & - & 0 & 0 & 0 & 0 & + & + & - & - & 0 & 0 \\ + & - & - & + & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & + & - & - & + & 0 & 0 \\ 0 & 0 & 0 & 0 & + & - & - & + & + & - & - & + & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}.$$



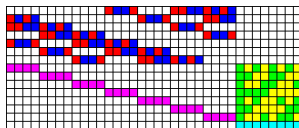
## Some Recent Progress on ETFs

[Jasper, Mixon, F 14] Every McFarland harmonic ETF is a rotated Steiner ETF. New infinite family of optimal codes.

[F, Mixon, Jasper 16] New infinite family of complex ETFs arising from finite projective planes containing hyperovals.

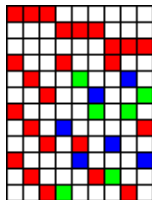


[F, Jasper, Mixon, Peterson 18]: Tremain's construction of an  $\text{ETF}(15, 36)$  generalizes. New infinite family of *real* ETFs.

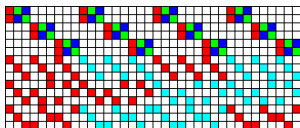


## Some More Recent Progress on ETFs

[F, Jasper, King, Mixon 18] Some ETFs can be represented in terms of the regular simplices they contain.



[F, Jasper 19] Generalizing Davis-Jedwab difference sets gives new infinite families of ETFs from group divisible designs.



# Some Future Directions

**Fundamental mysteries:** Lifting, spectral estimation.

**Some mature open problems:**

- ▶ Zauner's conjecture.
- ▶ Optimal projective packings when no ETF/OGF exists.
- ▶ Integrality conditions on the existence of complex ETFs.
- ▶ Breaking the square-root bottleneck for deterministic RIP.

**Not-so-high hanging fruit:**

- ▶ New constructions of ETFs, OGFs, ECTFFs, EITFFs.
- ▶ New connections to combinatorial designs.

# Literature

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Thank you, John! Happy Birthday!

